

情報論理学 II 試験 3

2003 年 1 月 31 日 (金)

1. CL 項 $\lambda^*xy.xyy$ を求めよ (S, K で表せ)。
2. CL 項 $(\lambda^*xy.xyy)xy$ を weak reduction せよ。
3. CL 項 $\lambda^*xy.xyy$ が持つ型 δ を求め, $\lambda^*xy.xyy \in \delta$ に至る TA_C の証明図をかけ。
4. ラムダ項 $\lambda xy.(\lambda uz.y(xz))(\lambda y.yz)$ が持つ型 δ を求め,
 $\lambda xy.(\lambda uz.y(xz))(\lambda y.yz) \in \delta$ に至る TA_λ の証明図をかけ。
5. ラムダ項 $\lambda xy.(\lambda uz.y(xz))(\lambda y.yz)$ の β 正規形 (β -nf) M を求め, 4 で求めた δ に対して $M \in \delta$ に至る TA_λ の証明図をかけ。

定義 0.1 (Weak reduction) Any CL-term KXY or $SXYZ$ is called a (weak) redex. Contracting an occurrence of a redex in a term U means replacing one occurrence of

$$\begin{array}{l} KXY \quad \text{by} \quad X, \\ SXYZ \quad \text{by} \quad XZ(YZ). \end{array}$$

If this changes U to U' , we say that U (weakly) contracts to U' , or

$$U \triangleright_{1w} U'.$$

We say that U (weakly) reduces to V , or

$$U \triangleright_w V$$

iff V is obtained from U by a finite (perhaps empty) series of weak contractions.

定義 0.2 (Abstraction) For each CL-term M and each variable x , a CL-term $\lambda^*x.M$ is defined by induction on M , thus:

- (a) $\lambda^*x.M \equiv KM$ if $x \notin FV(M)$;
- (b) $\lambda^*x.x \equiv SKK$;
- (c) $\lambda^*x.Ux \equiv U$ if $x \notin FV(U)$;
- (d) $\lambda^*x.UV \equiv S(\lambda^*x.U)(\lambda^*x.V)$ if neither (a) nor (c) applies.

定義 0.3 (The type-assignment system \mathbf{TA}_C) It has two axiom-schemes; they are

$$\begin{array}{l} (\rightarrow K) \quad K \in \alpha \rightarrow \beta \rightarrow \alpha, \\ (\rightarrow S) \quad S \in (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma. \end{array}$$

Its only rule is called the \rightarrow -elimination rule or ($\rightarrow e$);

$$\frac{X \in \alpha \rightarrow \beta \quad Y \in \alpha}{XY \in \beta} (\rightarrow e).$$

定義 0.4 (The type-assignment system \mathbf{TA}_λ) It has no axiom-schemes and two rules the \rightarrow -elimination rule ($\rightarrow e$) and the \rightarrow -introduction rule ($\rightarrow i$);

$$\frac{X \in \alpha \rightarrow \beta \quad Y \in \alpha}{XY \in \beta} (\rightarrow e),$$

$$\frac{[x \in \alpha] \quad M \in \beta}{\lambda x.M \in \alpha \rightarrow \beta} (\rightarrow i).$$